

#### Public transport revenue optimization under no-elongation and no-stopover restrictions (Results of the Thesis) Seet Viskari 13.6.2025

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Työn saa tallentaa ja julkistaa Aalto-yliopiston ayoimilla yerkkosiyuilla. Muilta osin kaikki oikeudet pidätetään





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#### **Background (1/5): Motivation**

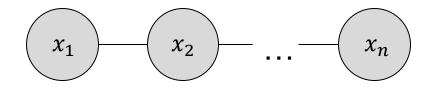
- Public transportation providers aim to maximize revenue
- Providers determine optimal prices for connections based on passengers' willingness-to-pay
- Special interest in no-elongation and no-stopover properties
- They produce a more consistent fare system and reduce the possibility for exploitation





#### **Background (2/5): PTN**

- Public transport network (PTN) modeled as a path graph (V, E)
- Set of stations  $V = \{x_1, x_2, \dots, x_n\}$
- Set of direct connections (paths) E







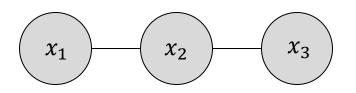
#### **Background (3/5): Ticket**

- Tickets for path  $W = [x_i, x_j]$ 
  - Standard ticket T(W)
  - Elongated ticket T(H), where  $W \subseteq H$
  - Compound ticket  $T(W_1, W_2)$ , where  $W_1 + W_2 = W$
- Price of a ticket  $T = (H_1, ..., H_t)$  denoted  $p_T = \sum_{j=1}^t p_{H_j}$



## Background (4/5): No-elongation constraint

 It cannot be cheaper for a passenger to buy a ticket for a longer path and use only a part of it.



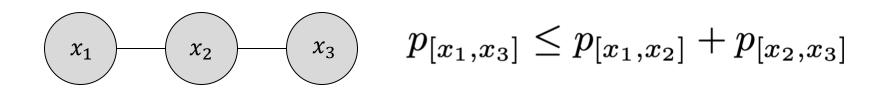
 $p_{[x_1,x_2]} \le p_{[x_1,x_3]}$ 





## Background (5/5): No-stopover constraint

 It cannot be cheaper for a passenger to split a path into multiple sections and buy separate tickets for each subpath









- Build a model that solves the optimal prices for tickets to maximize the provider's revenue
- Illustrate how the no-elongation and no-stopover properties affect the optimal solution





### Methodology (1/2): Assumptions

- Each path  $[x_i, x_j]$  has some demand  $s_{[x_i, x_j]}$  (number of passengers)
- Demand is distributed into groups with different valuations (willingess-to-pay) for the path
- Distributions
  - Equal (all groups have equal number of passengers)
  - Decreasing (groups with high valuations have less passengers)
  - Random (demand is randomly split into groups)



#### Methodology (2/2): Model

- Formulate models
  - CPM: no-elongation and no-stopover
  - UPM: prices have no constraints
  - EPM: passengers exploit the prices of UPM
- Compare the optimal solutions between different models
- Compare the optimal solutions between different distributions of valuations





#### Model (1/3): Constrained Pricing Model (CPM)

$$\max \sum_{\substack{x_i, x_j \in V \\ i \neq j}} \sum_{g=1}^G \Gamma_{[x_i, x_j]}^g \cdot s_{[x_i, x_j]}^g \quad (Maximize the revenue)$$
s.t. 
$$\sum_{\substack{x_i, x_j \in V \\ i \neq j}} \sum_{g=1}^G s_{[x_i, x_j]}^g \cdot \delta_{[x_i, x_j]}^g \ge B \quad (Minimum quota for tickets sold)$$

$$\Gamma_{[x_i, x_j]}^g \le p_{[x_i, x_j]} \qquad \forall x_i, x_j \in V, \ i \neq j, \ \forall g \in \{1, \dots, G\}$$

$$\Gamma_{[x_i, x_j]}^g \le M \cdot \delta_{[x_i, x_j]}^g \qquad \forall x_i, x_j \in V, \ i \neq j, \ \forall g \in \{1, \dots, G\}$$

$$p_{[x_i, x_j]} \le v_{[x_i, x_j]}^g + M \cdot (1 - \delta_{[x_i, x_j]}^g) \qquad \forall x_i, x_j \in V, \ i \neq j, \ \forall g \in \{1, \dots, G\}$$

$$p_{[x_i, x_j]} \le p_{[x_i, x_k]} + p_{[x_k, x_j]} \qquad \forall x_i, x_j, x_k, x_j \in V, \ i \neq j, \ \forall g \in \{1, \dots, G\}$$

$$p_{[x_i, x_j]} \le p_{[x_i, x_k]} + p_{[x_k, x_j]} \qquad \forall x_i, x_j, x_k, x_j \in V, \ i \neq j, \ \forall g \in \{1, \dots, G\}$$

$$p_{[x_i, x_j]} \le p_{[x_k, x_l]} \qquad \forall x_i, x_j, x_k, x_l \in V, \ k \le i < j \le l \text{ or } k \ge j \ge l$$

$$\delta_{[x_i, x_j]}^g \in \mathbb{R}_+ \qquad \forall x_i, x_j \in V, \ i \neq j, \ \forall g \in \{1, \dots, G\}$$

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## Model (2/3): Unconstrained Pricing Model (UPM)

 $\max \sum_{\substack{x_i, x_j \in V \\ z \neq z}} \sum_{g=1}^G \Gamma^g_{[x_i, x_j]} \cdot s^g_{[x_i, x_j]}$ (N-E) and (N-S) are removed s.t.  $\sum_{\substack{x_i, x_j \in V \\ i \neq j}} \sum_{g=1}^G s^g_{[x_i, x_j]} \cdot \delta^g_{[x_i, x_j]} \ge B$  $\forall x_i, x_j \in V, \ i \neq j, \ \forall g \in \{1, ..., G\}$ (12)  $\Gamma^g_{[x_i, x_j]} \le p_{[x_i, x_j]}$  $\forall x_i, x_j \in V, \ i \neq j, \ \forall g \in \{1, ..., G\}$  $\Gamma^{g}_{[x_i, x_i]} \leq M \cdot \delta^{g}_{[x_i, x_i]}$  $p_{[x_i, x_j]} \le v_{[x_i, x_j]}^g + M \cdot (1 - \delta_{[x_i, x_j]}^g) \quad \forall x_i, x_j \in V, \ i \neq j, \ \forall g \in \{1, ..., G\}$  $\delta^{g}_{[x_i, x_j]} \in \{0, 1\}$  $\forall x_i, x_j \in V, i \neq j, \forall q \in \{1, ..., G\}$  $\forall x_i, x_i \in V, i \neq j$  $p_{[x_i,x_j]} \in \mathbb{R}_+$  $\Gamma^g_{[x_i,x_i]} \in \mathbb{R}_+$  $\forall x_i, x_j \in V, \ i \neq j, \ \forall g \in \{1, ..., G\}$ 





# Model (3/3): Exploited Pricing Model (EPM)

- Passengers exploit the optimal ticket prices given by UPM
- Out of ticket prices given by UPM, manually find the cheapest ticket (standard, elongated, or compound) for each path
  - Elongated: consider all elongated tickets and choose the cheapest one
  - Compound: shortest path algorithm





## Results (1/2): Comparing optimal solutions

Distribution		_	Average	# Tickets	# Paid max
of groups	$\mathbf{Model}$	Revenue	price	$\mathbf{sold}$	valuation
Equal	$\operatorname{CPM}$	34444	30.5	1171	368
	UPM	38212	34	1169	385
	$\mathbf{EPM}$	26800	20	1299	385
Decreasing	$\operatorname{CPM}$	28299	28	1181	405
	UPM	31157	29.5	1170	407
	$\mathbf{EPM}$	26800	20	1299	433
Random	$\operatorname{CPM}$	41820	34.5	1293	368
	UPM	51470	44.5	1293	490
	EPM	38900	30.5	1350	287

Number of stations = 5 and number of groups = 5





#### **Results (2/2): Performance analysis**

Runtime (s)

0.06 0.15 0.18 0.22 0.22 0.29

0.12 0.26 0.36 0.58 0.46 0.98

0.22 0.69 1.04 1.90 1.50 2.84

► - 0.21 1.25 2.74 4.54 7.81 5.23

ω - 0.38 2.94 7.40 12.67 26.85 25.71

Groups

decreasing

m - 0.10 0.08 0.05 0.13 0.13 0.09

4 -

<u>ν</u> η -

5tat 6

- 1.2

- 1.0

0.8

- 0.6

- 0.4

- 0.2

0.0

3 4 5 6 - 25

- 20

- 15

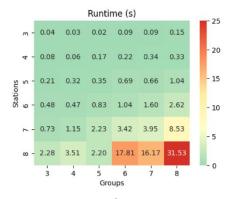
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- 5

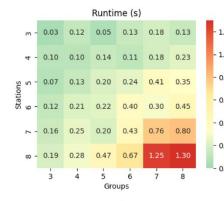
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0.0

7 8

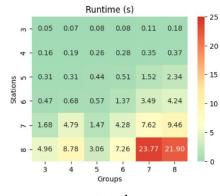


random



random





equal

Runtime (s)

0.05 0.05 0.09 0.12 0.11 0.11

0.10 0.18 0.25 0.23 0.36 0.39

0.16 0.26 0.41 0.76

0.29 0.56

Groups

► - 0.22 0.29 0.31 0.49 0.81

0.07 0.08 0.18 0.18 0.29

0.08

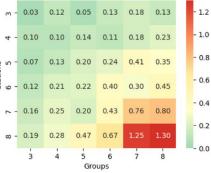
0.15

∞ - 0.26 0.40

3 4 5 6

4 -

Stations 6 5



#### UPM

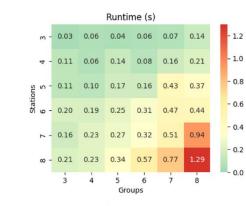


1.10 1.33

7 8







decreasing

#### Conclusions

- Introducing the no-elongation and no-stopover constraints yields higher revenue, assuming passengers would otherwise exploit the system
- Benefits the provider but passengers face higher prices and less tickets are sold
- Future study: complex graphs, bi-objective (e.g. revenue and tickets sold)







Anita Schöbel and Reena Urban. The cheapest ticket problem in public transport. Transportation Science, 56(6):1432–1451, 2022.



